

Research Question

Can we use **dynamic programming as a model-based paradigm** for combinatorial optimization?

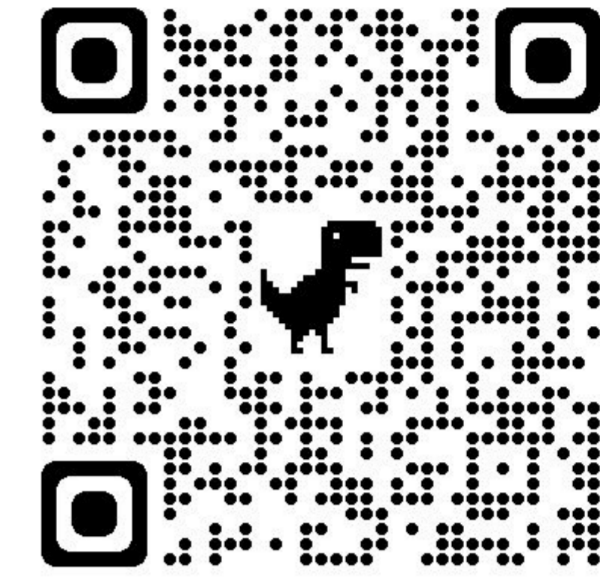
Developed Software

Use our software to **solve your problem by just defining a DP model!**

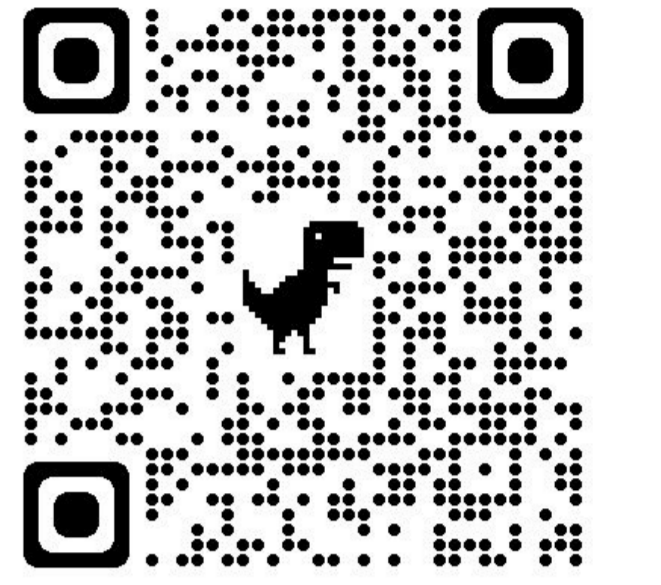
We developed general-purpose **heuristic search solvers!**

Install the Python interface: `pip install didppy`

Open-source and free for commercial use (MIT/Apache-2.0).



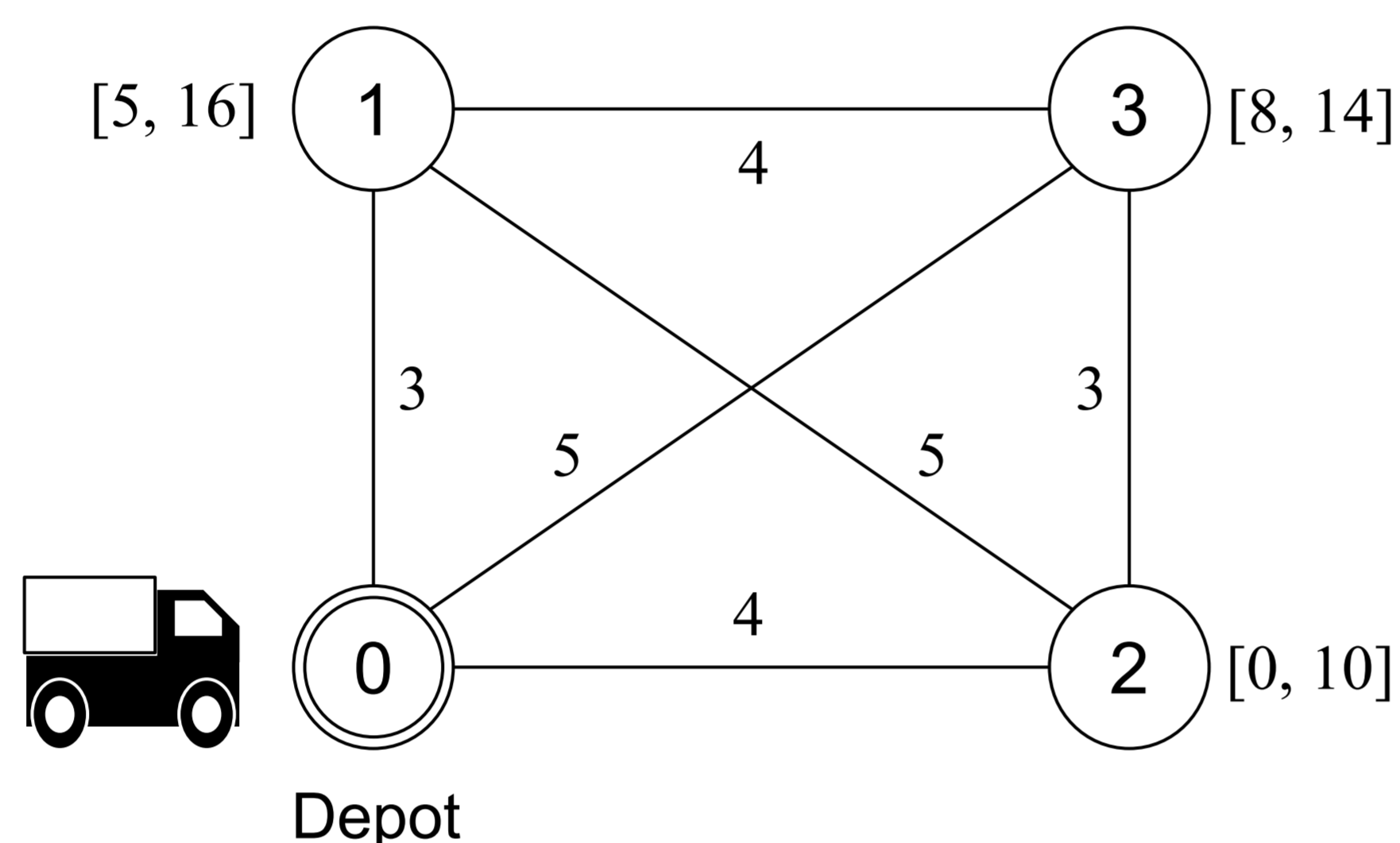
Project page



Tutorials, examples, and API reference

TSPTW Example

Minimize the travel time of a tour to visit all customers within the time windows.



Modeling and Solving in DIDP

```
import didppy as dp

model = dp.Model(maximize=False)

customer = model.add_object_type(number=4)
a = [0, 5, 0, 8]
b = [100, 16, 10, 14]
c = model.add_int_table([[0, 3, 4, 5], [3, 0, 5, 4], [4, 5, 0, 3], [5, 4, 3, 0]])

u = model.add_set_var(object_type=customer, target=[1, 2, 3])
i = model.add_element_var(object_type=customer, target=0)
t = model.add_int_resource_var(target=0, less_is_better=True)

for j in range(1, 4):
    visit = dp.Transition(
        name="visit {}".format(j),
        cost=c[i, j] + dp.IntExpr.state_cost(),
        effects=[(u, u.remove(j)), (i, j), (t, dp.max(t + c[i, j], a[j]))],
        preconditions=[u.contains(j), t + c[i, j] <= b[j]],
    )
    model.add_transition(visit)

return_to_depot = dp.Transition(
    name="return",
    cost=c[i, 0] + dp.IntExpr.state_cost(),
    effects=[(i, 0), (t, t + c[i, 0])],
    preconditions=[u.is_empty(), i != 0],
)
model.add_transition(return_to_depot)

model.add_base_case([u.is_empty(), i == 0], cost=0)

for j in range(1, 4):
    model.add_state_constr(~u.contains(j) | (t + c[i, j] <= b[j]))

model.add_dual_bound(0)

solver = dp.CABS(model)
solution = solver.search()
```

DP Model for TSPTW

compute $V(N \setminus \{0\}, 0, 0)$

$$V(U, i, t) = \begin{cases} \min_{j \in U: t + c_{ij} \leq b_j} c_{ij} + V(U \setminus \{j\}, j, \max\{t + c_{ij}, a_j\}) & \text{if } U \neq \emptyset \\ c_{i0} + V(\emptyset, 0, t + c_{i0}) & \text{else if } i \neq 0 \\ 0 & \text{otherwise} \end{cases}$$

State variables:

- U : unvisited customers
- i : current customer
- t : current time

Constants

- N : all customers (0: depot)
- $[a_i, b_i]$: time window for customer i
- c_{ij} : travel time from customer i to j

What DIDP Can Do but PDDL Cannot

Explicitly modeling implications of the problem definition that can be useful solvers (common in OR!).

State constraints (can be used for pruning)

$$V(U, i, t) = \infty \text{ if } \exists j \in U, t + c_{ij} > b_j$$

```
for j in range(1, 4):
    model.add_state_constr(~u.contains(j) | (t + c[i, j] <= b[j]))
```

Dominance with resource variables (can be used for pruning)

$$V(U, i, t) \leq V(U, i, t') \text{ if } t \leq t'$$

```
t = model.add_int_resource_var(target=0, less_is_better=True)
```

Dual bound (can be used as a heuristic)

$$V(U, i, t) \geq 0$$

```
model.add_dual_bound(0)
```

Heuristic Search Solvers

Heuristic search solves a DP model as a shortest path problem in a **state space** using the dual bound as a heuristic.

We developed the following solvers:

- CAASDy: A*.
- CABS: performs beam search with exponentially increasing beam width (anytime and complete).
- 5 other anytime heuristic search solvers.

Promising performance compared to MIP and CP.

Future work: parallelization, domain-independent dual bound

	Description	MIP	CP	CAASDy	CABS
TSPTW (340)	TSP with time	227/0.227	47/0.026	257/0.244	259/0.003
CVRP (207)	vehicle routing	26/0.585	0/0.317	5/0.976	6/ 0.185
SALBP-1 (2100)	assembly line	1357/0.345	1584/0.005	1653/0.213	1801/0.000
Bin Packing (1615)	bin packing	1157/0.039	1234/0.002	922/0.429	1163/ 0.002
MOSP (570)	manufacturing	225/0.039	437/0.004	483/0.153	527/0.000
Graph-Clear (135)	building security	24/0.110	4/0.015	76/0.437	103/0.000
Talent Scheduling (1000)	scheduling actors	6/0.051	7/ 0.002	224/0.793	253/0.011
m-PDTSP (1117)	pick up & delivery	945/0.078	1049/0.013	947/0.196	1035/ 0.002
$1 \sum w_i T_i$ (375)	job scheduling	109/0.018	150/ 0.000	270/0.280	285/0.034

Coverage / primal gap (gap to the best known cost) achieved within 8GB and 30-min.